Lecture 22

Wednesday Nov. 29

Program:

```
inc_by_9
  require  i > 3
  do       i > 5
    i := i + 9
  ensure
    i > 13
  end
```

Is this correct?

4. $i > 5 \Rightarrow i > 4$ ☐ (No e.g. $i = 4$; $i > 3$)

existing ✓ (Proved) wp ☐

1. Translate prog. into Hoare Triple

$\{ i > 5 \}$ $i := i + 9$ $\{ i > 13 \}$ ($i > 3$)

2. Prove existing precond. $i > 5$ is no weaker than $i > 3$

$wp(i := i + 9, i > 13)$

3. Calculate $wp(i := i + 9, i > 13)$

$wp(i := i + 9, i > 13)$ (post-state)

$= \{$ wp rule for assign. $\}$

$i > 13 \; [i := i_0 + 9]$

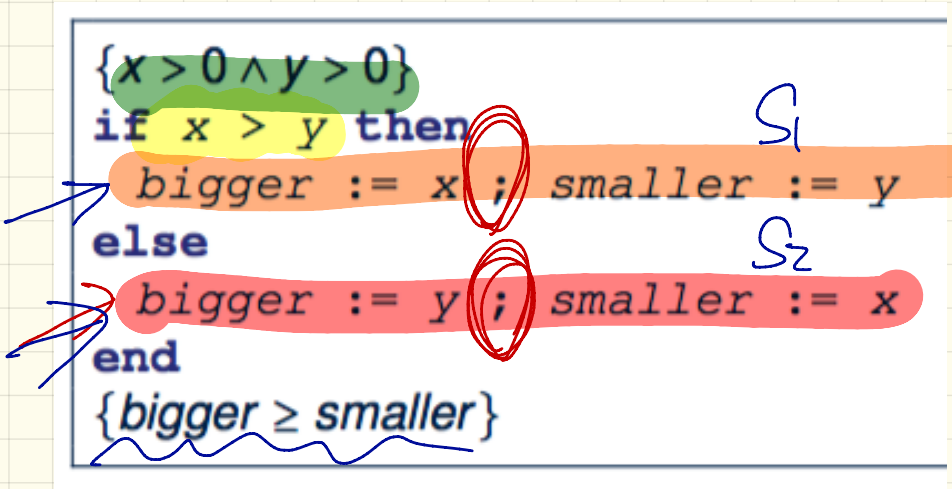
post-state value

$= \{$ subs. $\}$

$i_0 + 9 > 13$

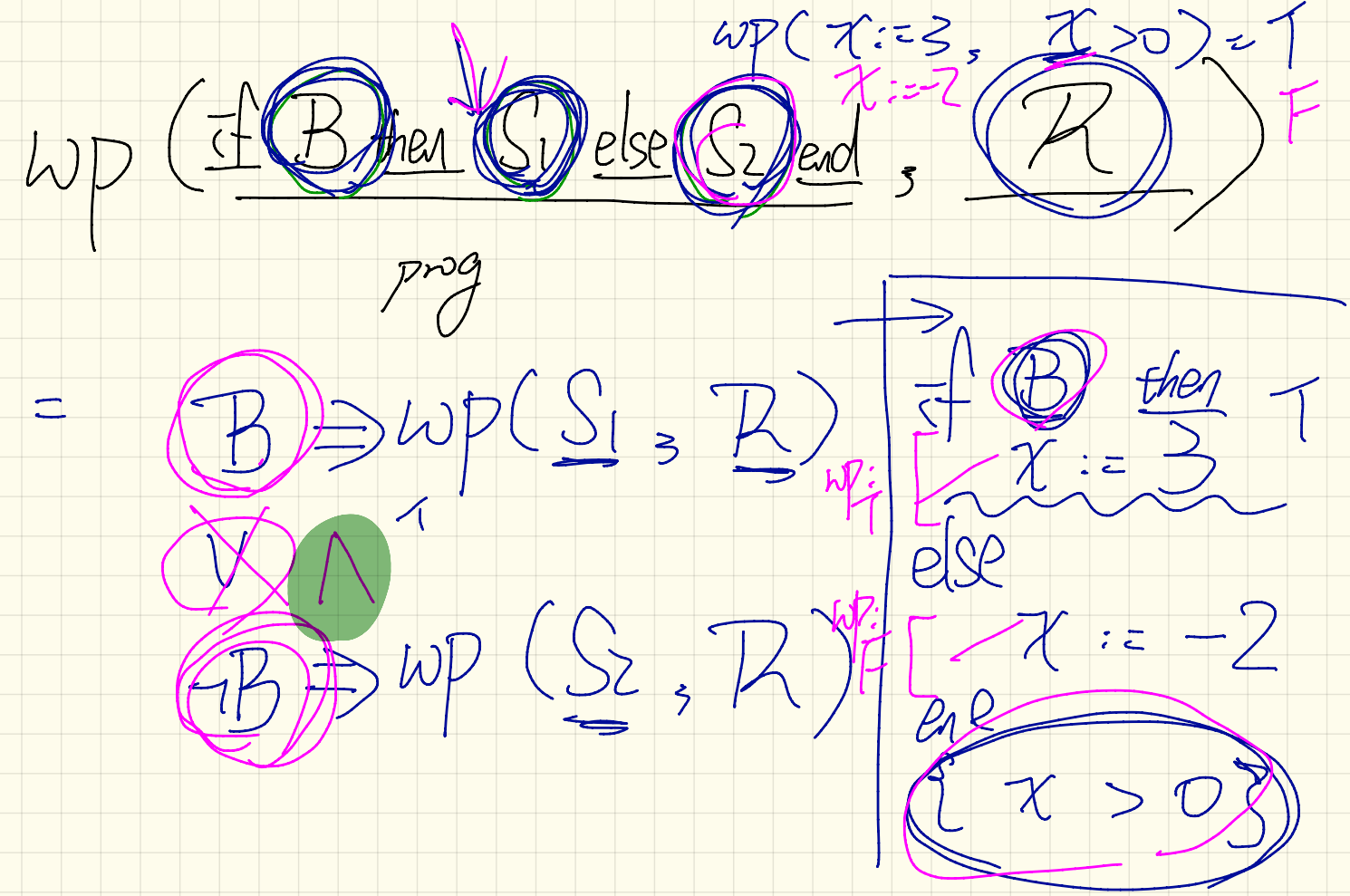
$= \{$ simp. $\}$

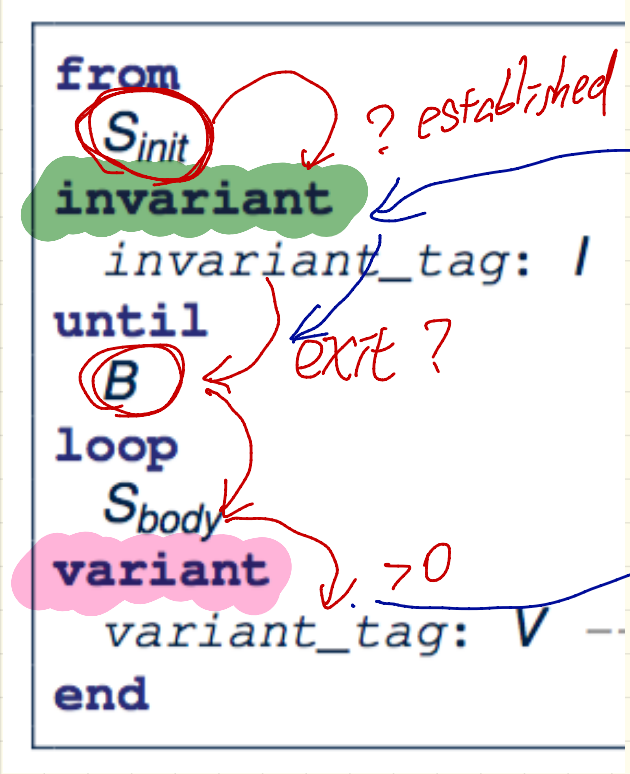
$i_0 > 4$ ← wp for $i := i + 9$ to establish $i > 13$

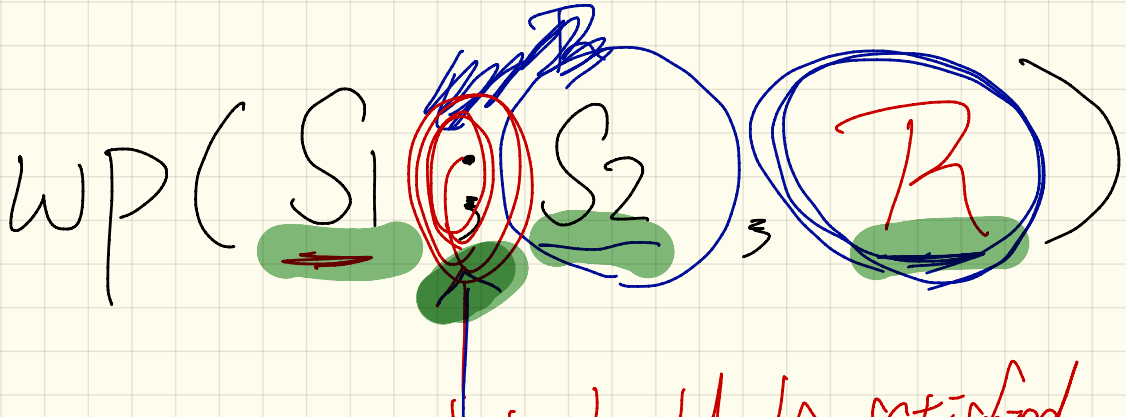


$\{x>0 \wedge y>0 \wedge x>y\}\ S_1\ \{bigger \geq smaller\}$

$\{x>0 \wedge y>0 \wedge \neg(x>y)\}\ S_2\ \{bigger \geq smaller\}$


wp(x := 3, x > 0) = T
x := -2
F
wp (if B then S1 else S2 end ; R)
prog
=
B ⇒ wp(S1 ; R)
∧
¬B ⇒ wp (S2 ; R)
wp;
wp;
if B then T
x := 3
else
x := -2
end
{ x > 0 }


from
S_{init}
invariant
invariant_tag: I
until
B
loop
S_{body}
variant
variant_tag: V
end
? established
? maintained
exit ?
>0


wp(S1 ; S2 , R)
what should be satisfied
(starting precondition for S2)
established by S1
wp(S1 ; wp(S2 , R))
starting pointing of S2
ending point of S1

$$wp(S_1 ; S_2, R)$$

$$wp(S_1, wp(S_2, R))$$

starting condition for S_2 to establish R.

$\{True\}\ tmp := x;\ x := y;\ y := tmp\ \{x > y\}$

1. $wp(tmp := x;\ x := y;\ y := tmp,\ x > y)$

$= \{$ wp for seq. comp. $\}$

$wp(tmp := x,\ wp(x := y;\ y := tmp,\ x > y))$

$= \{$ wp for seq. comp. $\}$

$wp(tmp := x,\ wp(x := y,\ wp(y := tmp,\ x > y)))$

$= \{$ wp for assign. $\}$

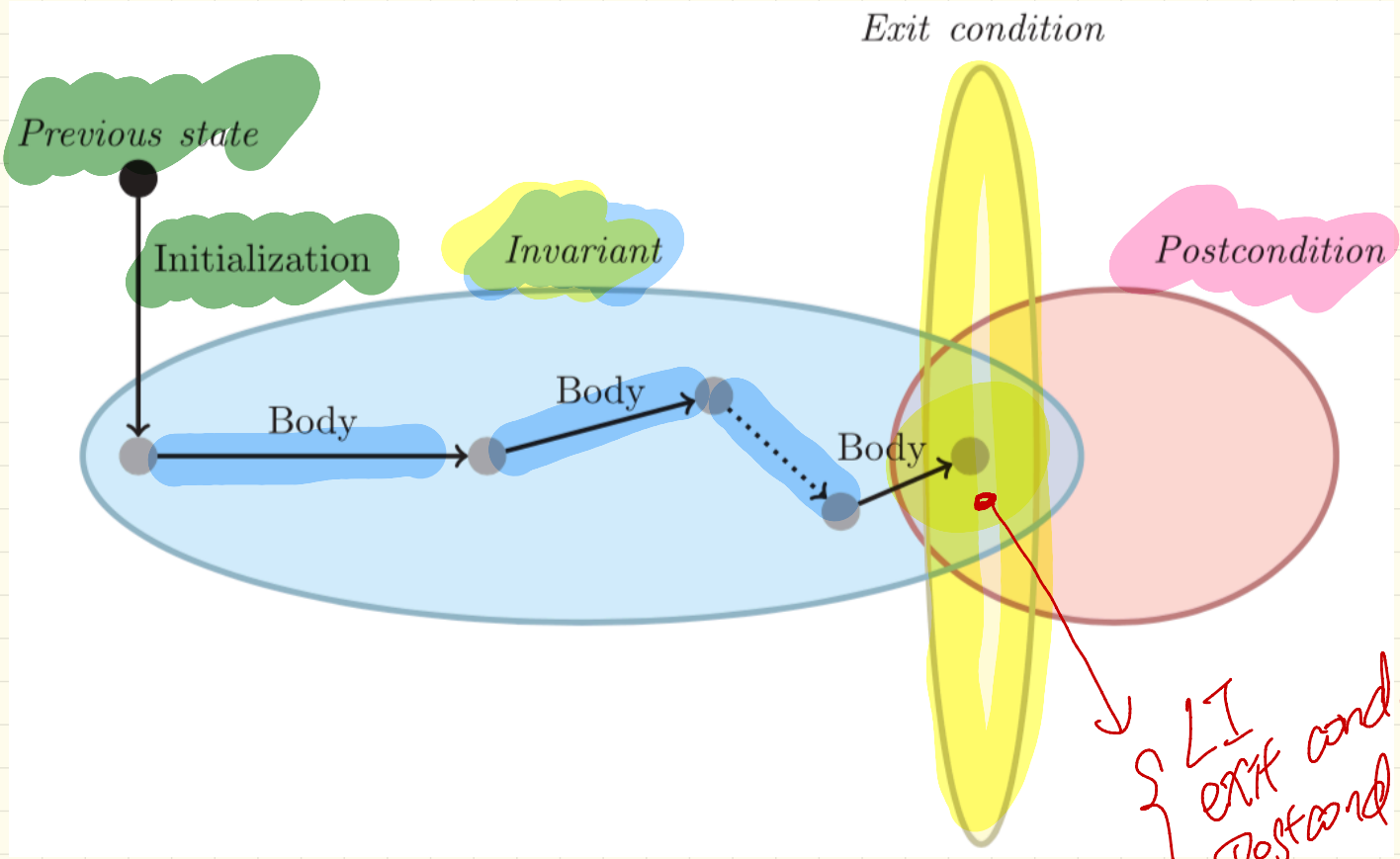
$wp(tmp := x,\ wp(x := y,\ x > y\ [y := tmp]))$

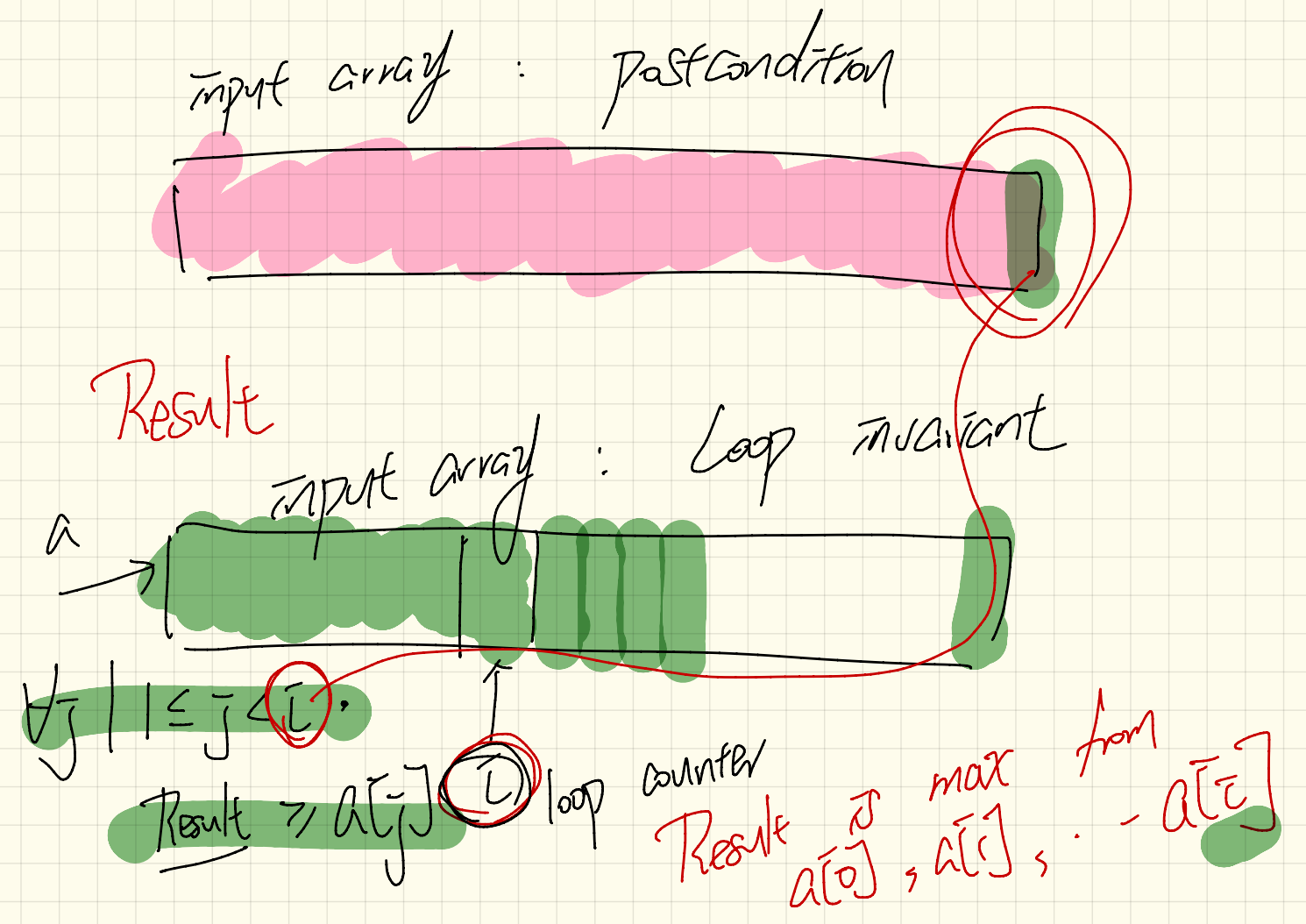
$= \{$ wp for assign $\}$

$wp(tmp := x,\ y > tmp) = y > x$

$x \geq tmp$

$True \Rightarrow y \geq x$


Exit condition
Previous state
Initialization
Invariant
Postcondition
Body
Body
Body
LI
exit cond.
postcond.


input array : postcondition
Result
input array : Loop invariant
a
∀j | 1 ≤ j < i ⇒
Result ≥ a[j]
i
loop counter
Result
max
from
a[0] , a[1] , . . . a[i]